



Accelerating Data Warehousing Applications Using General Purpose GPUs

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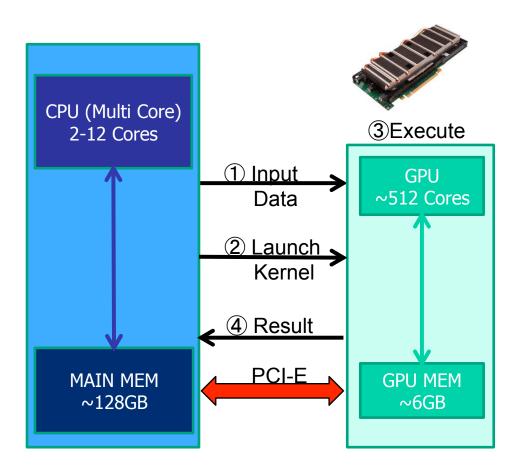
Sponsors: National Science Foundation, LogicBlox Inc., IBM, and NVIDIA

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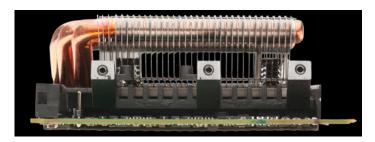
The General Purpose GPU

- GPU is a many core co-processor
 - 10s to 100s of cores
 - 1000s to 10,000s of concurrent threads
 - CUDA and OpenCL are the dominant programming models
- Well suited for data parallel apps
 - Molecular Dynamics, Options Pricing, Ray Tracing, etc.
- Commodity: led by NVIDIA, AMD, and Intel





Enterprise: Amazon EC2 GPU Instance



NVIDIA Tesla



Amazon EC2 GPU Instances

| Elements | Characteristics |
|----------|---|
| OS | CentOS 5.5 |
| CPU | 2 x Intel Xeon X5570 (quad-core "Nehalem" arch, 2.93GHz) |
| GPU | 2 x NVIDIA Tesla "Fermi" M2050 GPU Nvidia GPU driver and CUDA toolkit 3.1 |
| Memory | 22 GB |
| Storage | 1690 GB |
| I/O | 10 GigE |
| Price | \$2.10/hour |

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Data Warehousing Applications on GPUs

■The good

- Lots of potential data parallelism
- If data fits in GPU mem, 2x—27x speedup has been shown

■The bad

- Very large data set (will not even fit in host memory)
- I/O bound (GPU has no disk)
- PCI data transfer takes 15–90% of the total time*

| Order | Price | Discount |
|-------|-------|----------|
| 0 | 10 | 10% |
| 1 | 20 | 20% |
| 2 | 10 | 15% |
| 3 | 51 | 14% |
| 4 | 33 | 13% |
| 5 | 22 | 10% |
| | | |

* B. He, M. Lu, K. Yang, R. Fang, N. K. Govindaraju, Q. Luo, and P. V. Sander. Relational query co-processing on graphics processors. In TODS, 2009.



This Work

Goal: Enable Large data warehousing applications on GPUs

Assumptions

- In-memory system
 - Host memory, not GPU memory
- Not OLTP (*Online Transaction Processing*) type simple queries

Focus on data analysis instead of data entry/retrieval

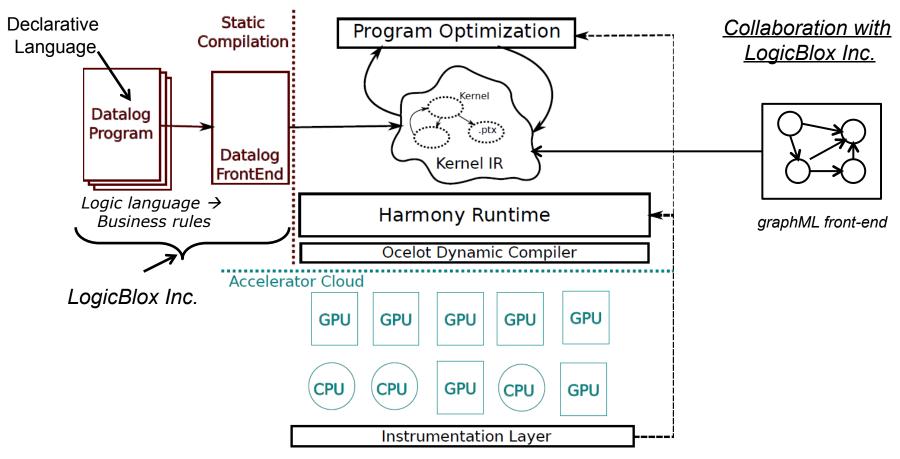


Research Thrusts

- I: Optimized implementations of primitives
 - Relational algebra (RA)
 - Data management within the GPU memory hierarchy
- II: Data movement optimizations
 - Between host and accelerators
 - Within an accelerator
- III: In-core processing
 - Cluster wide memory aggregation techniques
 - Change the ratio of host memory size to accelerator memory size



Red Fox: Execution Environment for the Enterprise



- Bridge the x86-based Database Enterprise platform and Database backend with NVIDIA accelerators
- 10x-100x targeted improvement in application speedup

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Thrust I: Optimized Primitives

- Optimized implementation of each relational algebra (RA) operator
 - Synthesized from micro-primitives
 - Implemented as a CUDA/PTX kernel template and available as a library
- The Redfox compiler synthesizes an application by instantiating templated skeletons of these primitives
 - Provides a framework for optimizations (e.g. kernel fusion)



Relational Algebra Operators in GPU

| Operator | NVIDIA C2050 | Phenom 9570 | Speedup |
|---------------|-----------------|----------------|---------|
| Inner join | 26.4-32.3 GB/s | 0.11-0.63 GB/s | > 42x |
| Select | 104.2 GB/s | 2.55 GB/s | 41x |
| Set operators | 45.8 GB/s | 0.72 GB/s | 64x |
| Projection | 54.3 GB/s | 2.34 GB/s | 23x |
| Cross product | 98.8 GB/s | 2.67 GB/s | 37x |

- 10 Datalog microbenchmarks running
- Metrics based on random data sets, compressed rows and 16M tuple relations
- Cost of initial sort not included



Status

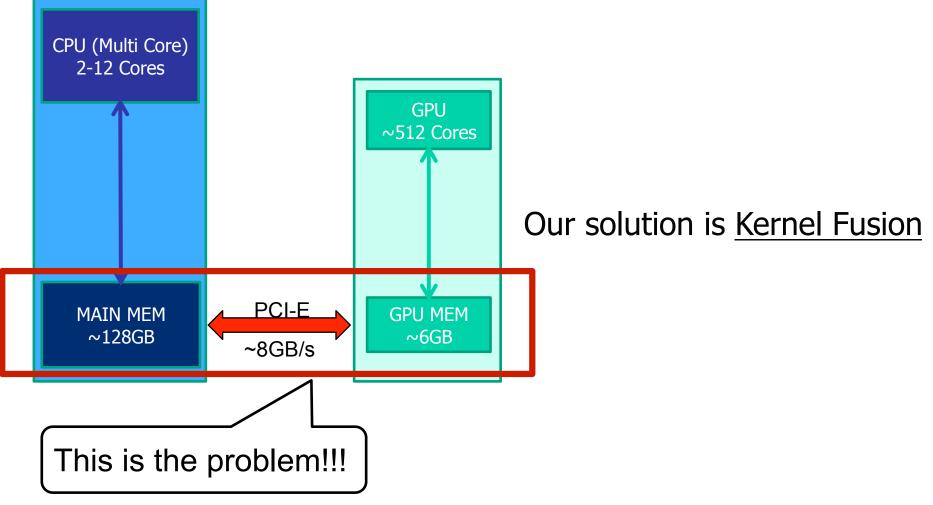
- Moving Red Fox to the Amazon EC2
- Robustness extensions across
 - Scale and size of tables
 - Size and diversity of data types
- Performance extensions
 - Single node and multi-node implementations
 - Ocelot remote device interface
 - Using multiple GPU configurations





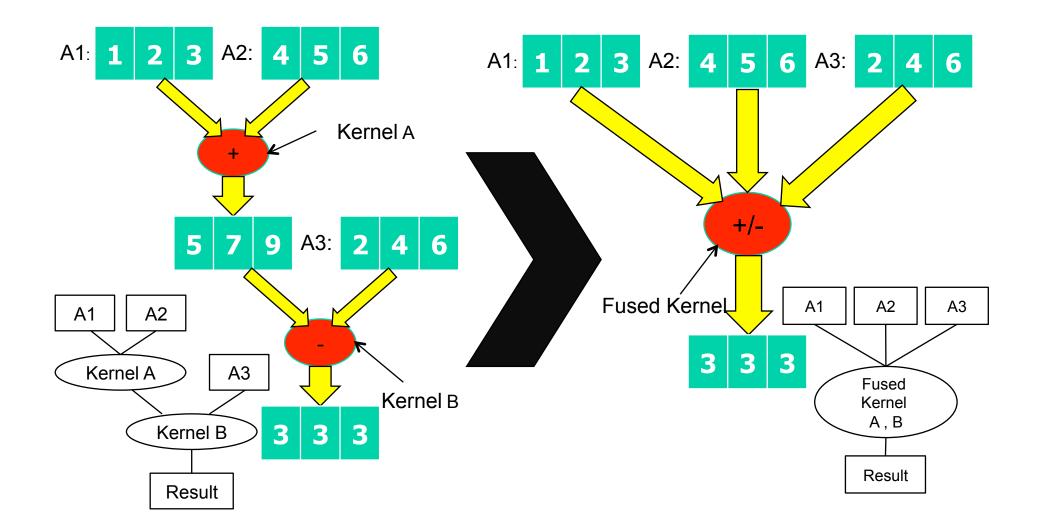
Thrust II: Optimization of Data Movement

<u>Collaboration with</u> <u>NEC Inc.</u>





Kernel Fusion

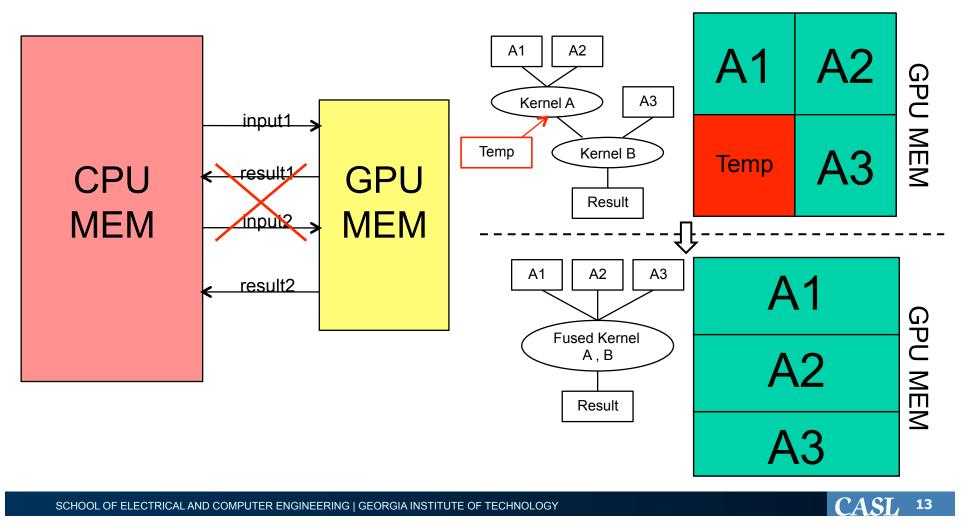




Benefits of Kernel Fusion-1

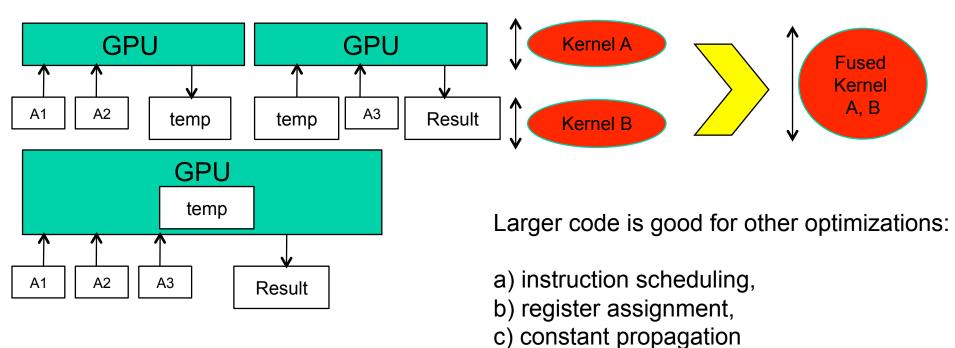
Reduce Data Transfer

Reduce Temp Storage



Benefits of Kernel Fusion-2

Faster ComputationEnable More Optimization

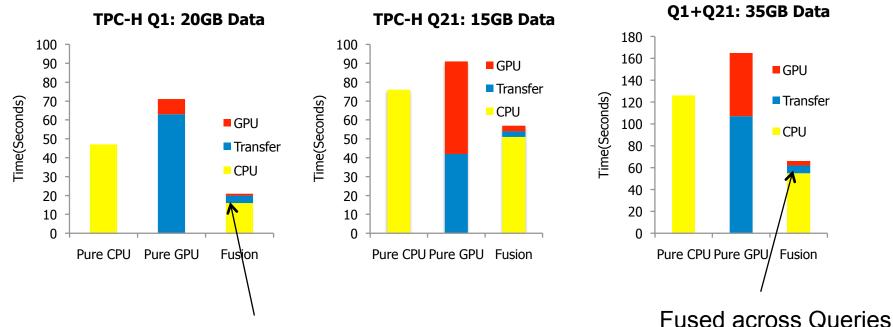


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Traverse the data only ONCE



Preliminary Result (2 Quad-Core CPU, C2070 GPU)

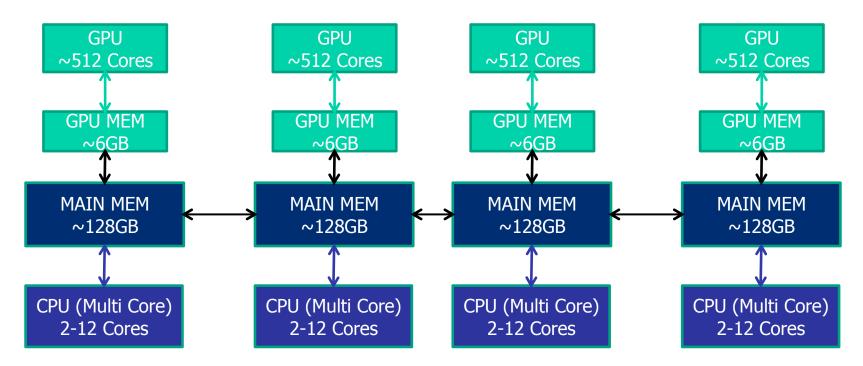


- Part of the query is run on CPU
- Transfer and GPU Computation time is much smaller





Thrust III: Cluster-based Memory Aggregation



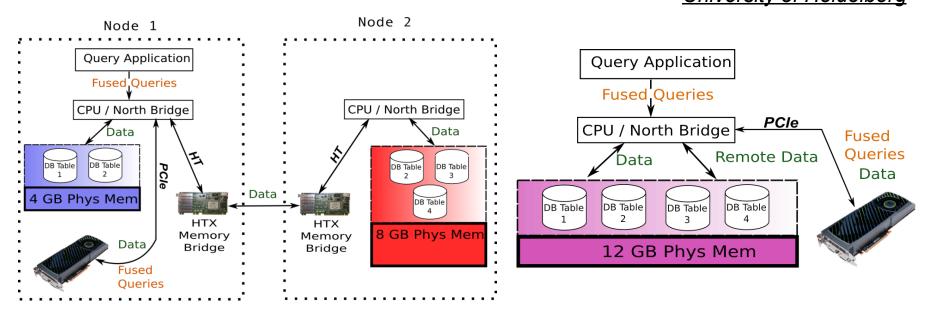
- Hardware support for global non-coherent, physical address space system
- Change the ratio of *host-memory* : GPU-memory



Global Address Space Support for In-Core Databases

<u>Collaboration with AIC Inc. &</u> University of Heidelberg

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- Use of low-latency, commodity network (HyperTransport) allows global, noncoherent access to remote memory
- Query app sees one large database / host memory from the application level
- Global address support can be extended in the future to support GPU memory
 - Applications could remotely read/write a remote GPU's memory without needing to involve its OS or CPU

Conclusions

- Fast GPU implementations of RA operators provide opportunity to run large data warehousing applications on GPU.
- Data movement optimization (Kernel Fusion) saves the memory transfer time and speeds up the computation time.
- New Memory Hierarchy (GAS) offers a larger logical memory for GPU database system.



Thank You

Questions?

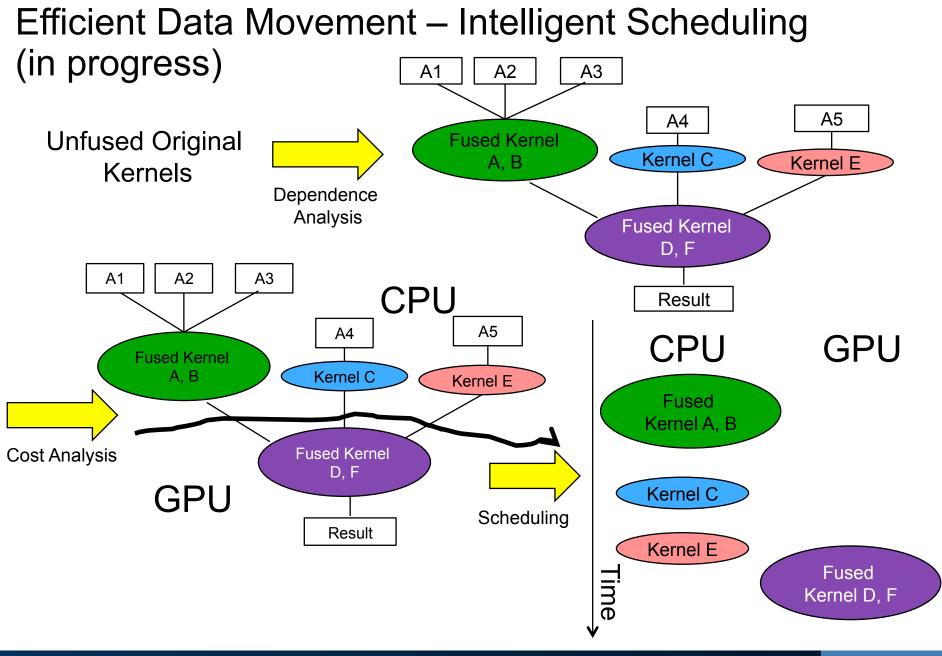


CASL ¹⁹

Backup

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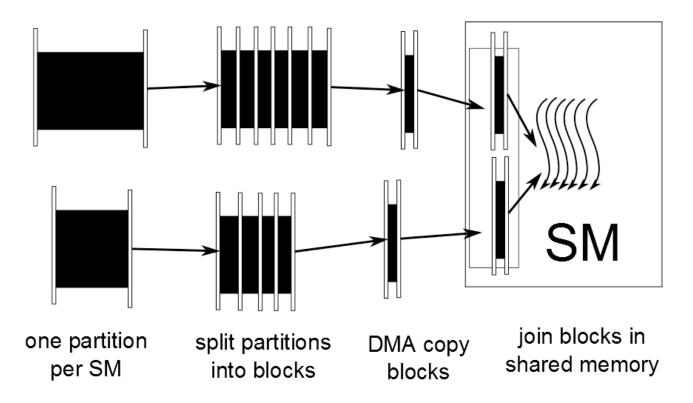




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CASL 21

Inner Join



Blocking into pages, shared memory buers, and transaction sized chunks makes memory accesses ecient.

