



# Putting Trust in Malicious Systems

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# Issues

**Defense:** can security software reliably use virtual environments?

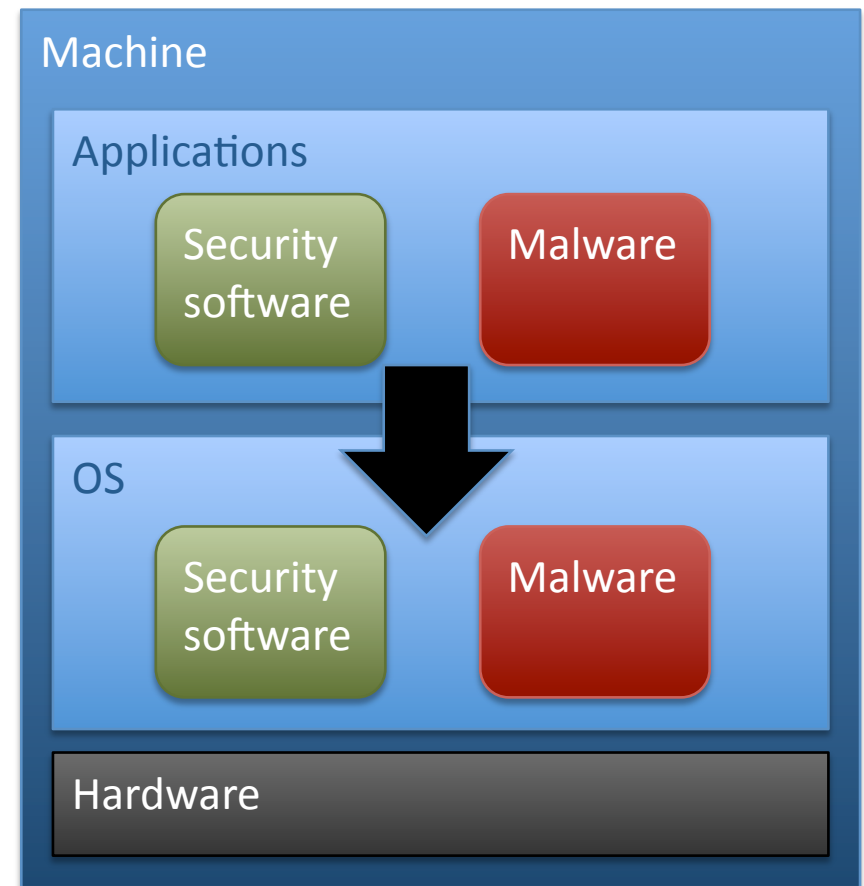
**Design:** what will hypervisors look like in five years?



# Traditional trust

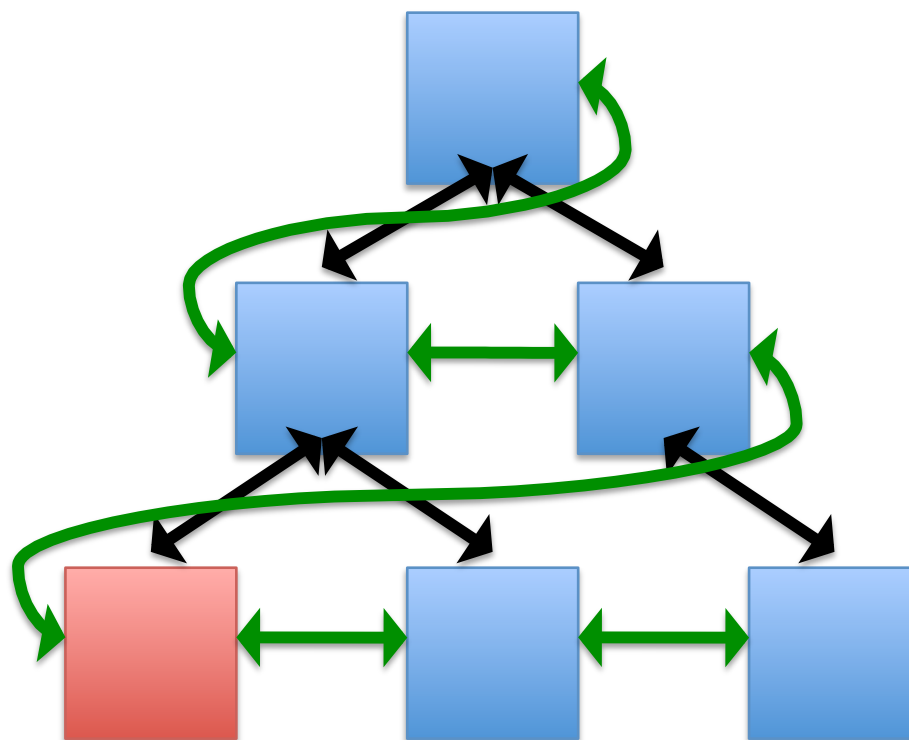
Truism: Whoever controls the lowest layer wins

Traditional software security architectures rely on unsafe trust relationship



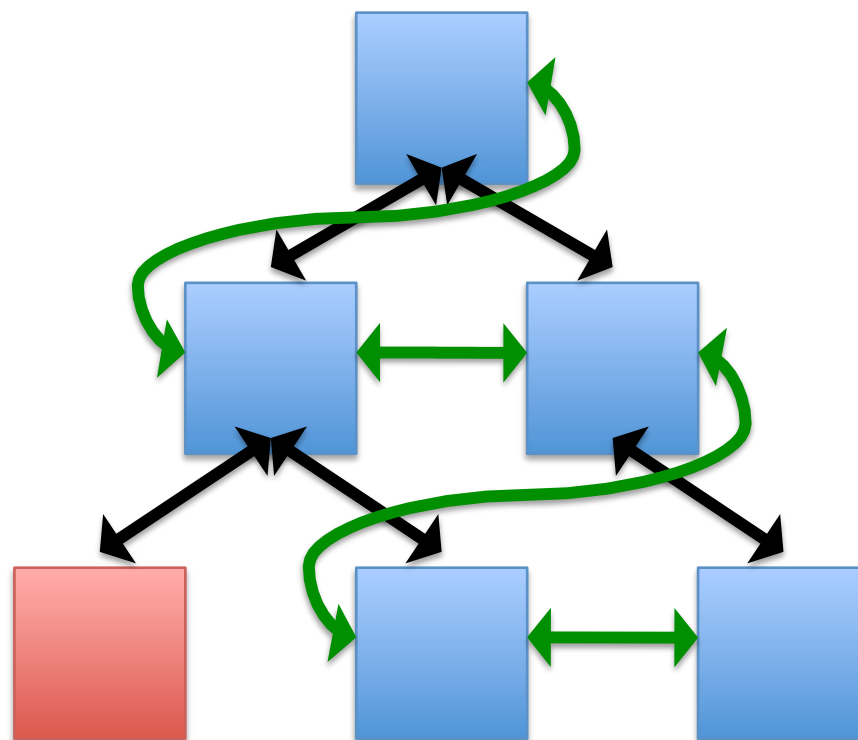
# Kernel-level malware

Unhook process from accounting list

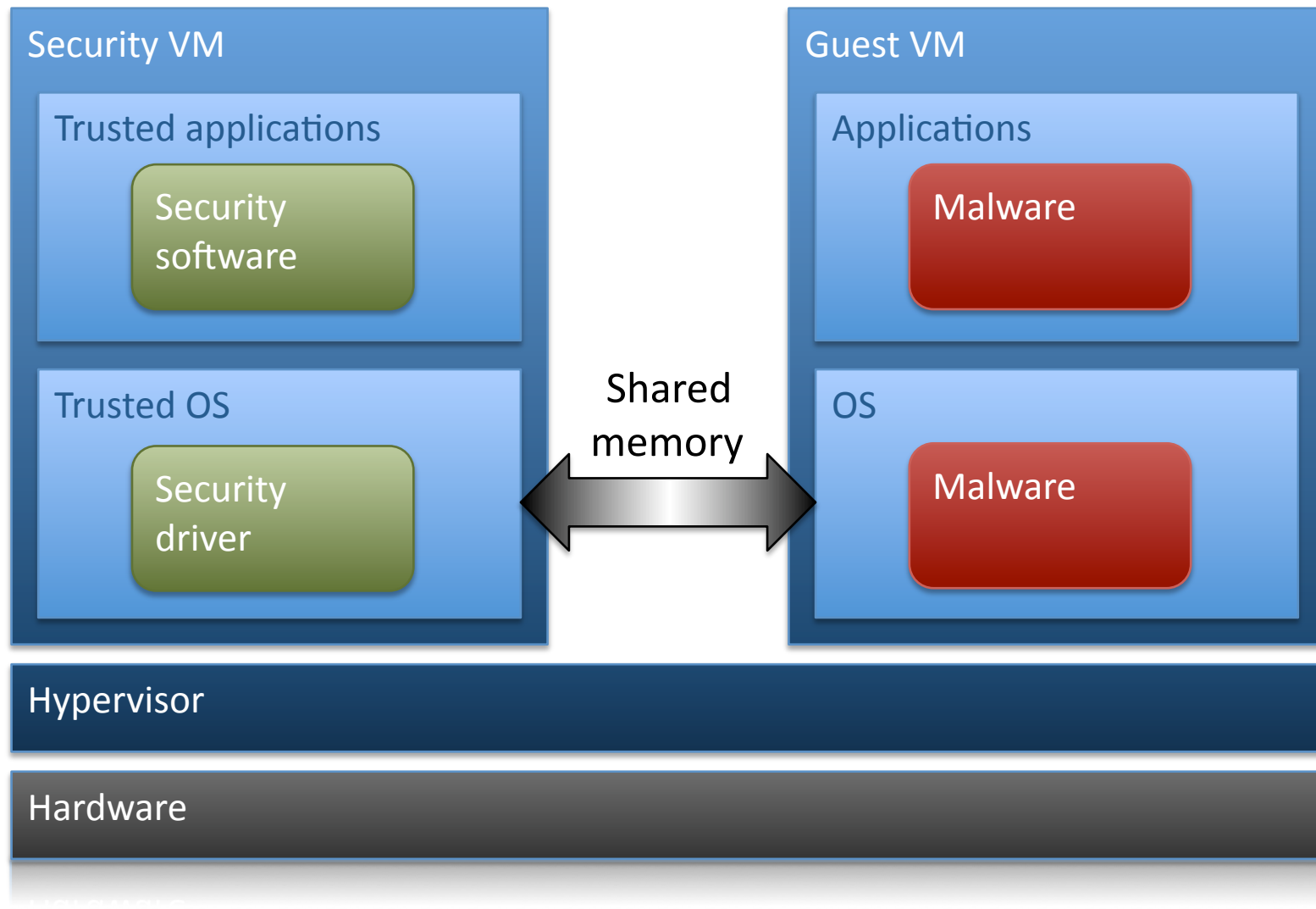


# Kernel-level malware

Unhook process from accounting list



# Modern architecture



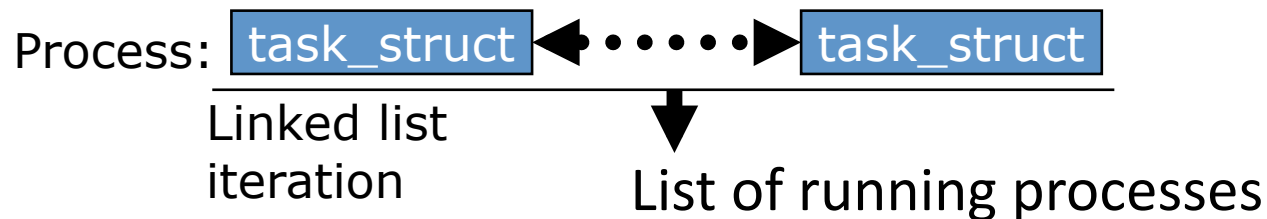
# Virtual machine introspection

0001010110101110101001011001100101110001110101110001110101000111010111100101000011010  
0110001000110100011101110101101100100010101101011101010010110011001011100011101011100  
0111010100011101011110010100001101001100010001101000111011101011011001000101011010111  
0101001011001100101110001110101110001110101000111010111100101000011010011000100011010  
0011101110101101100100010101101011101010010110011001011100011101011100011101010001110  
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1100100010101101011101010010110011001011100011101011100011101010001110101111001010000  
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0101100110010111000111010111000111010100011101011110010100001101001100010001101000111



# Virtual machine introspection

- Assign semantic meaning to raw memory bytes
  - Exported symbols
  - Predefined knowledge of kernel data types





# Virtual machine introspection

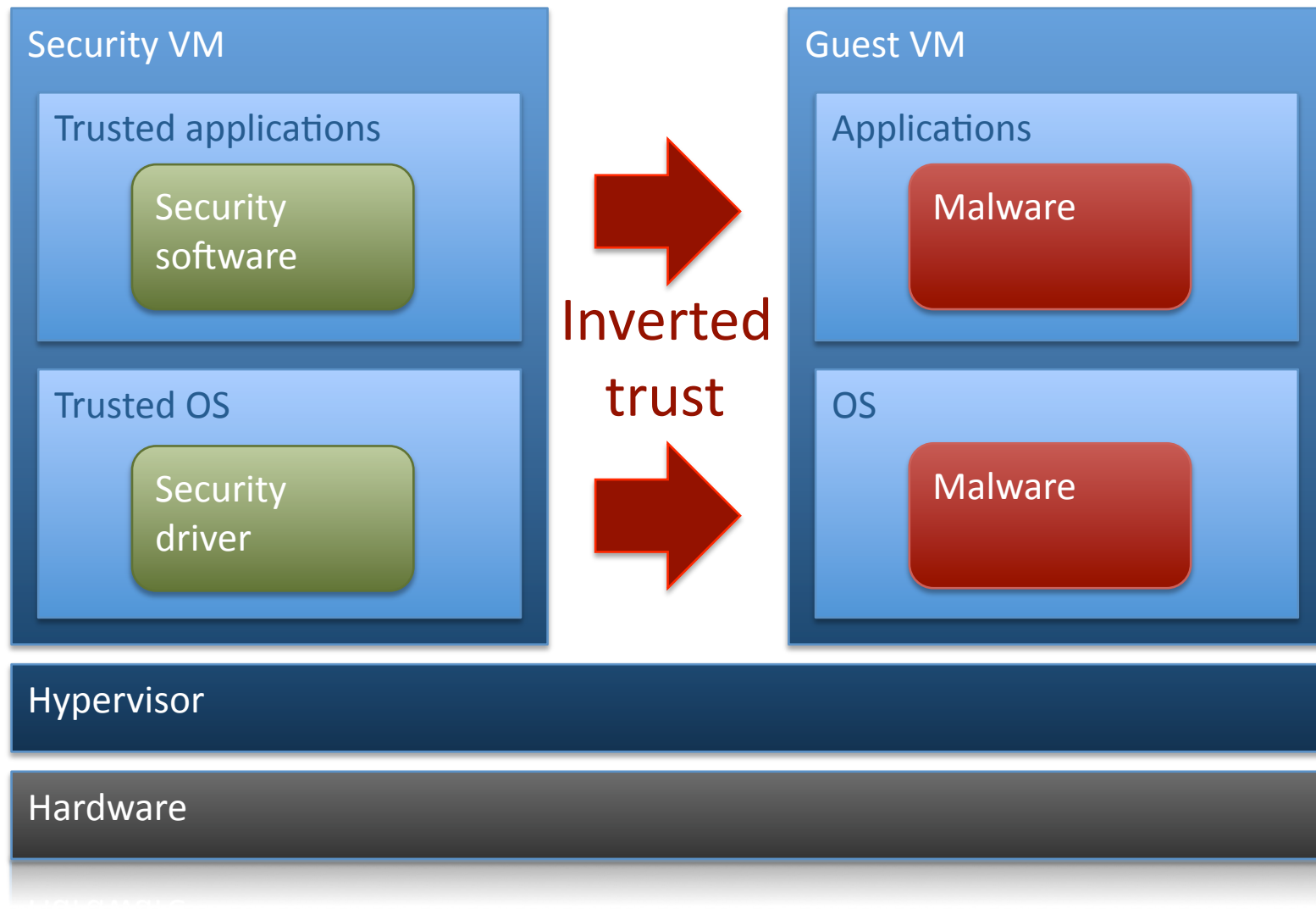
So... attackers can alter kernel  
dynamic data structures...

... and VMI-based utilities will build  
the incorrect view of victim VM...

... in the same way they would have had they  
been executing directly on the victim machine.



# Trust inversion

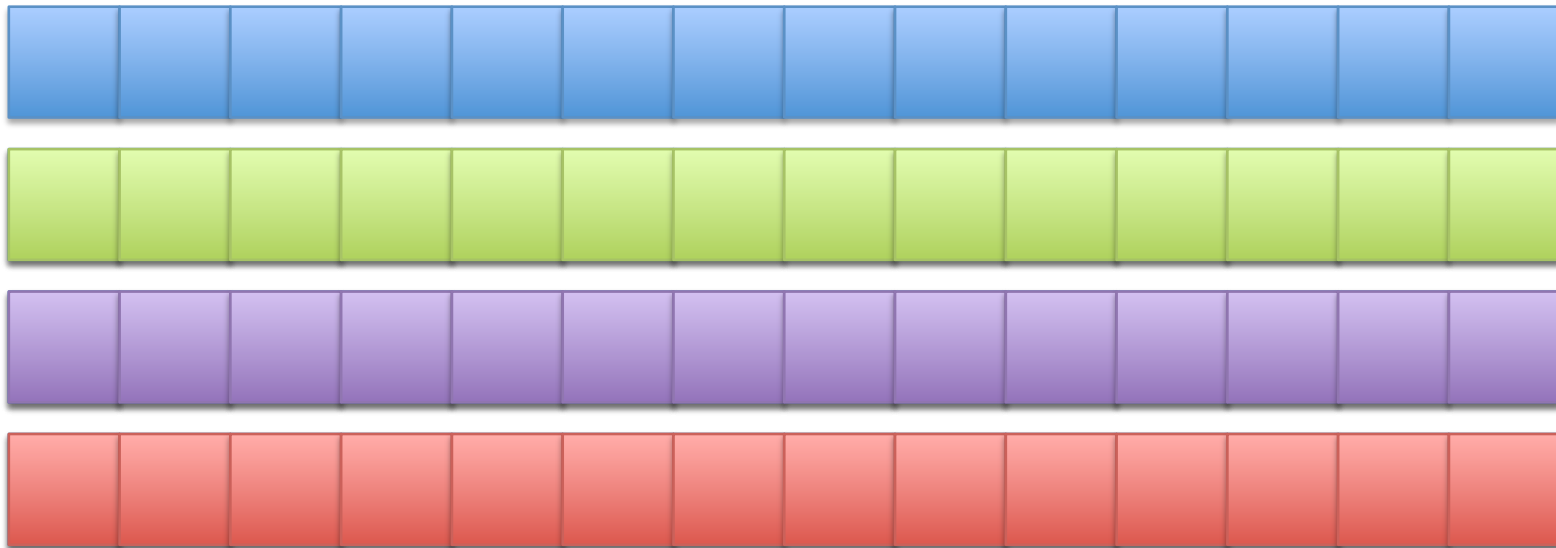


# Memory protection

Kernel memory:

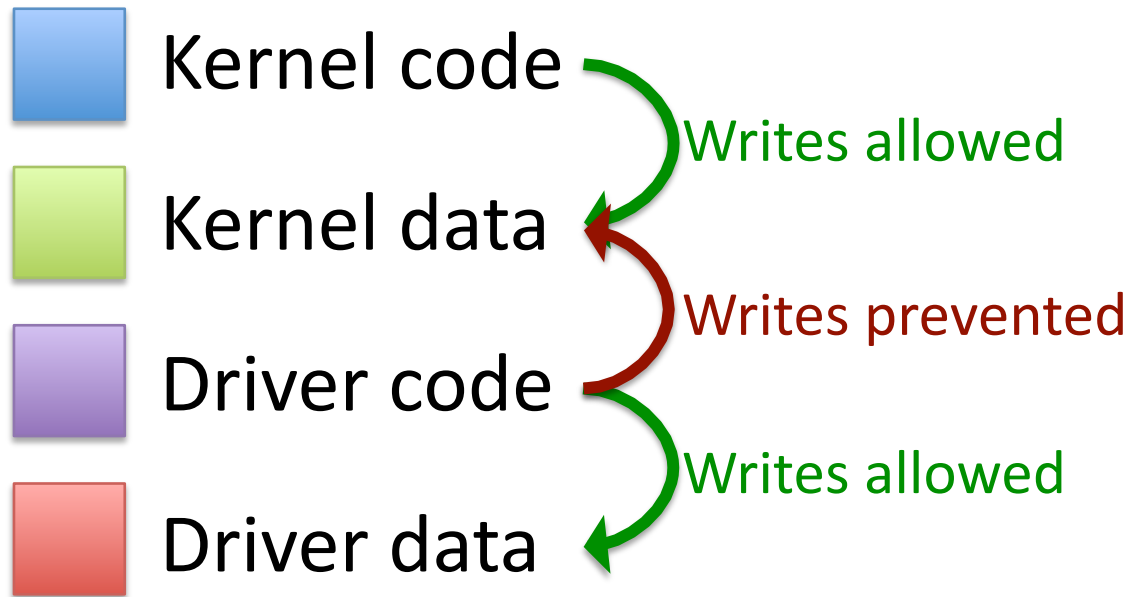


Process memory:



# Memory protection

Kernel memory:



# Kernel memory access control

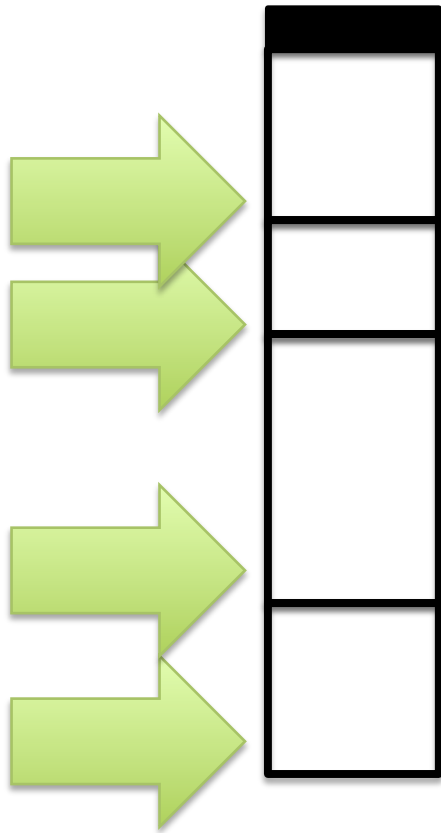
Kernel memory:



- Hypervisor & security VM enforce access control policy on kernel memory
- Policy: security-critical kernel data can be written only by core (non-driver) kernel code



# Policy enforcement



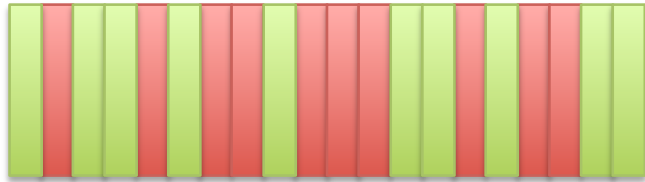
- How is provenance of write determined?
  - Stackwalk (pray for frame pointers)
  - Other techniques for corner cases



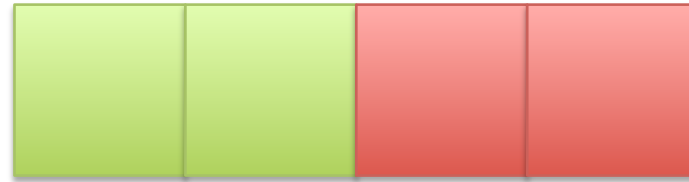
# Policy enforcement

- Reducing performance cost
  - Kernel data structure layout optimizations
  - Kernel memory allocation optimizations

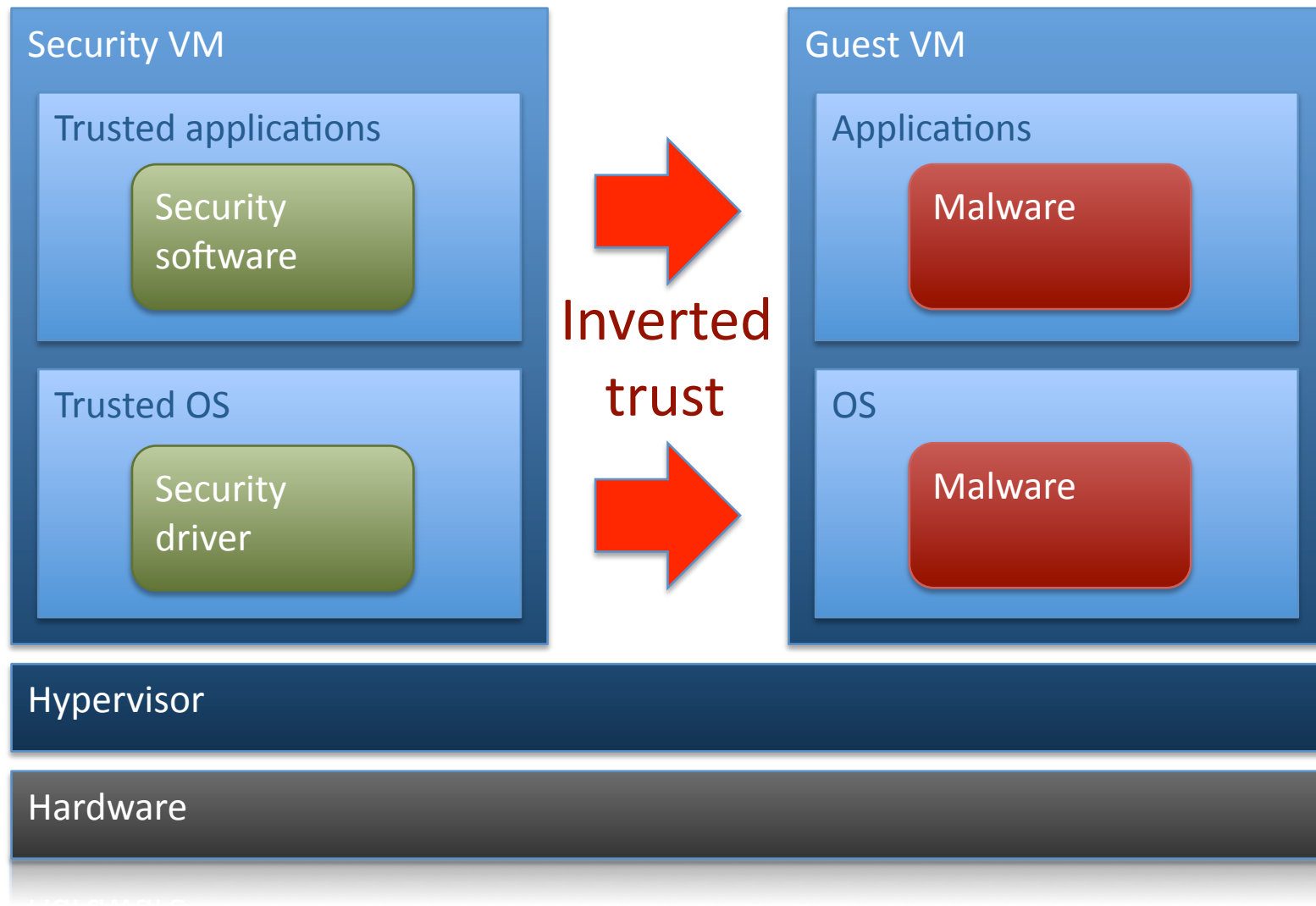
Make this:



Look like this:



# Kernel memory access control





# Conclusions

Trust assumptions made by VMI-based software require consideration

Hypervisors becoming kernels...  
kernels becoming libraries

Kernel-style memory protection  
restores trust in guest memory views





# Questions?

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